

Successive Refinement with Decoder Cooperation Himanshu Asnani, Haim Permuter, Tsachy Weissman



Stanford University, Ben-Gurion University, Stanford University



$\begin{array}{l} \textbf{Strictly-causal Cribbing}\\ \textbf{Rate Region}\\ \mathcal{R}_{sc}(D_1, D_2) = & \begin{cases} R_0 \geq \{I(X; \hat{X}_1, U) - H(\hat{X}_1 | U)\}^+ \\ R_0 + R_1 \geq I(X; \hat{X}_1, U) \\ \text{p.m.f.} : P(X, \hat{X}_1, U) \mathbf{1}_{\{\hat{X}_2 = f(U)\}} \\ \text{s.t. } \mathbf{E}[d_i(X, \hat{X}_i)] \leq D_i, \ i = 1, 2 \end{cases} \end{array}$

Achievability : "Forward Encoding and Block Markov Decoding"



• Source estimate of a decoder uses the estimate of another decoder, "crib".

• Non-causal, strictly-causal, causal cribbing.

Non-causal Cribbing

$$\mathcal{R}ate \ Region \\ \mathcal{R}_{nc}(D_1, D_2) = -\begin{cases} R_0 \ge \{I(X; \hat{X}_1, \hat{X}_2) - H(\hat{X}_1)\}^+ \\ R_0 + R_1 \ge I(X; \hat{X}_1, \hat{X}_2) \\ \text{p.m.f.} : P(X, \hat{X}_1, \hat{X}_2) \\ \text{s.t.} \ \mathbf{E}[d_i(X, \hat{X}_i)] \le D_i, \ i = 1, 2 \end{cases}$$



• Decoding at Decoder 1 requires $R_0 + R_1 \ge I(X; \hat{X}_1, \hat{X}_2)$



Extensions

- The results can be extended to the setting where "crib" is a deterministic function of the source reconstruction.
- Cooperation via "conferencing" or rate-limited link.

0.8

Theory, ISIT 2012.

• Dual channel coding setup : Multiple Access Channel (MAC) with common message and "*cooperating*" encoders.

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• Decoder 2 finds the unique bin "m" in which "crib" lies $\longrightarrow H(\hat{X}_1) \ge I(X; \hat{X}_1, \hat{X}_2) - R_0$

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